

Organic Self-organizing Bus-Based Communication Systems (OrganicBus)

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Hardware-Software-Co-Design

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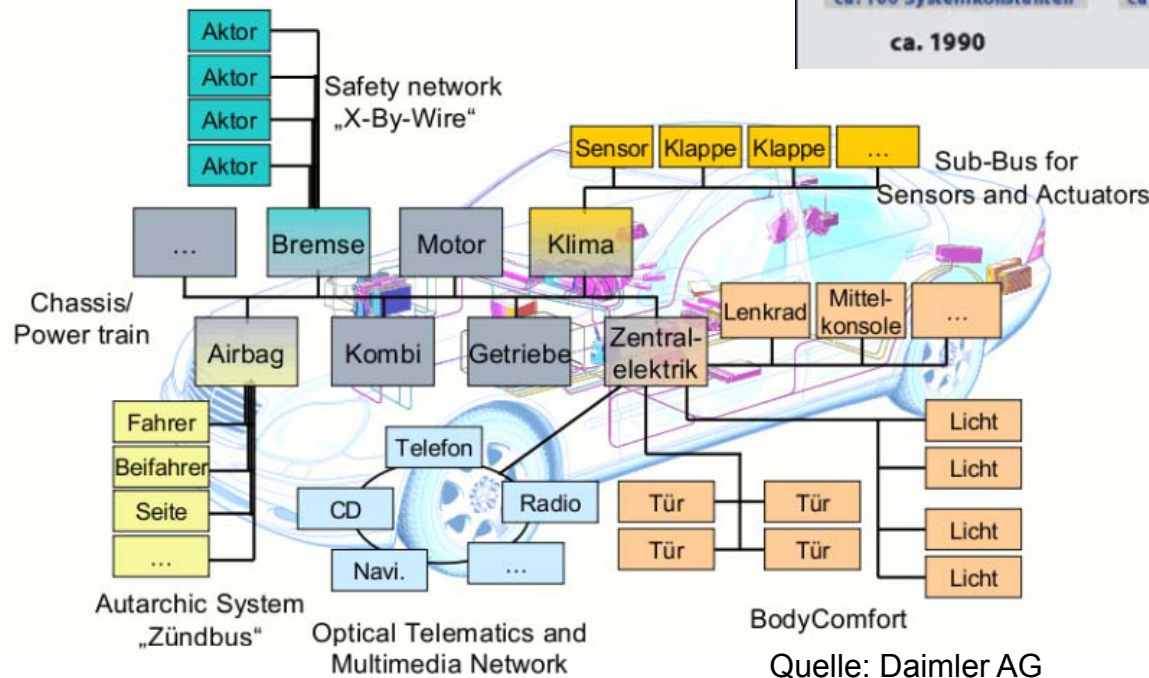
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Motivation and Idea

- Increasing complexity in distributed embedded systems
- Static bus arbitration too restricted or suboptimal



Quelle: VW



Quelle: Daimler AG

- Potential for optimization through *self-organizing bus systems (OrganicBus)*

Agenda

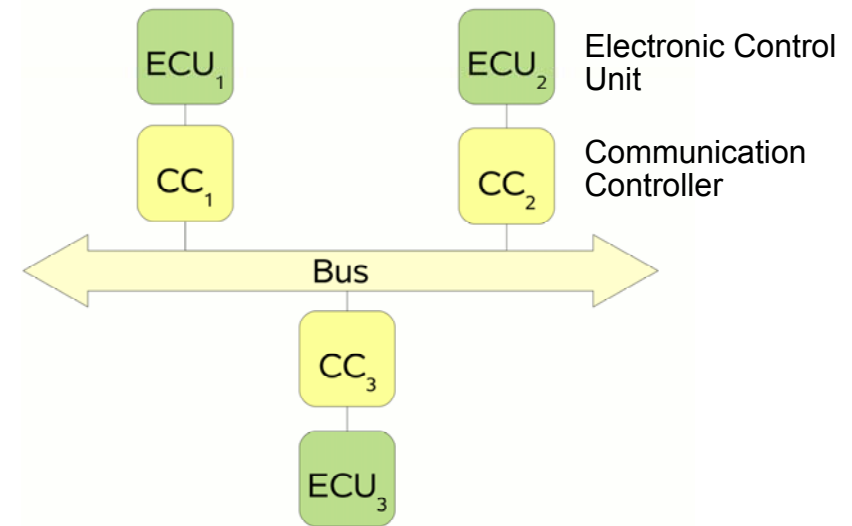
- Goals
- State of the art
- Roadmap
- First Milestones
- Conclusion

Goals

➤ Organic approach for bus-based communication systems

➤ Desired properties:

- Decentralized arbitration
- Self-organization
- Self-healing
- Self-optimization of bus access and traffic



➤ Separation of function and communication

➤ Reach the above properties using **simple local rules**

➤ Overcome the drawbacks of pure offline optimization

Major Goal: Increase overall Quality

- Typical types of messages to distinguish:
 - Hard deadline
 - Soft deadline
 - Bandwidth

- Definition of Quality (Objectives - in order of relevance):
 - Satisfaction of safety critical requirements
 - Increase of number of fulfilled constraints
 - Improvement of bus utilization
 - Guarantee of fairness

Features of our Decentralized Approach

- Apply multi-agent reinforcement learning techniques
- Apply game theoretic analysis for validation of protocols
- Use observer/controller architecture to preserve real-time properties
- Simulation and real hardware implementation
- Advantages:
 - Approach flexible to system changes
 - Robust to single node failures
 - Scalable

State of the Art

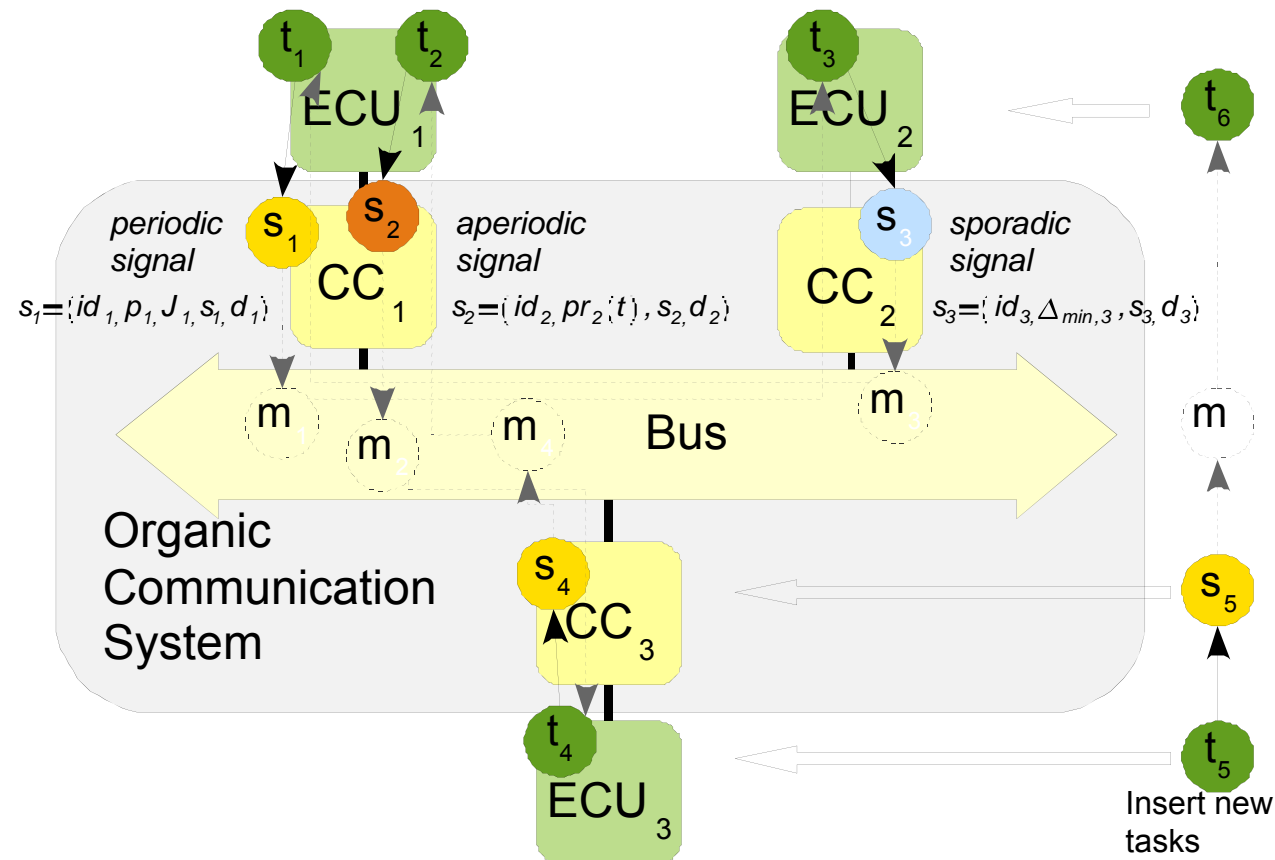
- Internet uses simple decentralized mechanisms, e.g. TCP
 - Fulfill requirements
 - Difficult to predict
- Embedded performance analysis for organic computing (EPOC) [Stein, Ernst]
 - Starting from offline solution
 - Hard real-time
- Feedback scheduling
 - Control loops optimize schedule
 - Centralized scheduling mechanism

Roadmap

1. Definition of a system model
2. Game-theoretic analysis
3. Development and rating of algorithms
4. Physical simulation
5. Implementation on real hardware

Definition of a System Model

- Separation between application and communication
- Describe messages by properties



Game-Theoretic Analysis

- Game Theory as a formal tool to define and analyze multi-agent decision problems
- A **game** consists of
 - a set of **players** $N = \{1, \dots, n\}$
 - a set of **strategies** S_i available to each player i
 - a specification of **payoffs** for each combination of strategies
- Payoff of player i : $u_i(\mathbf{s})$
 - strategies taken by all n players: $\mathbf{s} = (s_1, \dots, s_n)$
 - strategy of player i : $s_i \in S_i$

Equilibria

- *Nash Equilibrium* is an important characteristic to evaluate stable states in games
- When all players are selfish, a Nash Equilibrium is reached.
- *Selfish players* try to maximize their own payoff.
- No player achieves a higher payoff when choosing another strategy than his current strategy, when the other players keep their current strategies

Medium Access as 2-Player Game

- Models contention phase of medium access where users contend for medium access.
- Players have the strategies $\{send, wait\}$
- The payoff is
 - 1, if granted access to the medium
 - 0, else

		Player 2	
		wait	send
Player 1	wait	0,0	0,1
	send	1,0	0,0

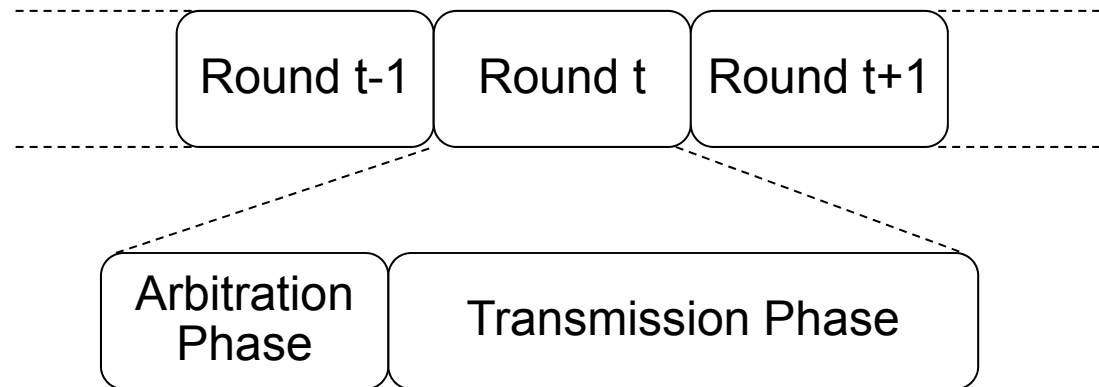
Contention-based access (e.g., WLAN),
Analysis in [1]

		Player 2	
		wait	send
Player 1	wait	0,0	0,1
	send	1,0	1,0

Priority-based Medium Access with
Player 1 having higher priority

Priority-based Access Game

- Arbitration and transmission phase define the rounds of the priority-based Access Game



- Mixed Strategies:
 - Not playing with discrete strategies “*wait*” and “*send*”.
 - Strategies are given as a probability distribution.
- Player i sends with probability p_i and waits with $1 - p_i$
- Payoff $u_i(\mathbf{s})$: Percentage of successfully sent messages

Unfairness of Selfish Users

- Goal is to achieve a *fair bandwidth sharing*, i.e.,

$$u_1(\mathbf{p}) = u_2(\mathbf{p}) = \dots = u_n(\mathbf{p})$$

- However, selfish behavior leads to unfair bandwidth sharing and blocking of users with low priorities.
- **Example:** 2-Player case

		Player 2	
		wait	send
Player 1	wait	0,0	0,1
	send	1,0	1,0

Nash Equilibria

Strategy “send” dominates strategy “wait” for Player 1

→ Player 1 always chooses “send”

→ Player 2 never sends

Enhancement to reach Fairness

- Main problems:
 - Users with low priorities cannot influence users with higher priorities.
 - Users with high priorities do not realize when blocking an user with lower priority.

- Idea: Demand that a small amount ε of the available bandwidth stays free.

- Payoff $u_i(\mathbf{s})$:
 - If sum of sending probabilities is less than $1 - \varepsilon$ then percentage of successfully sent messages
 - Else 0

Analysis of the Nash Equilibrium

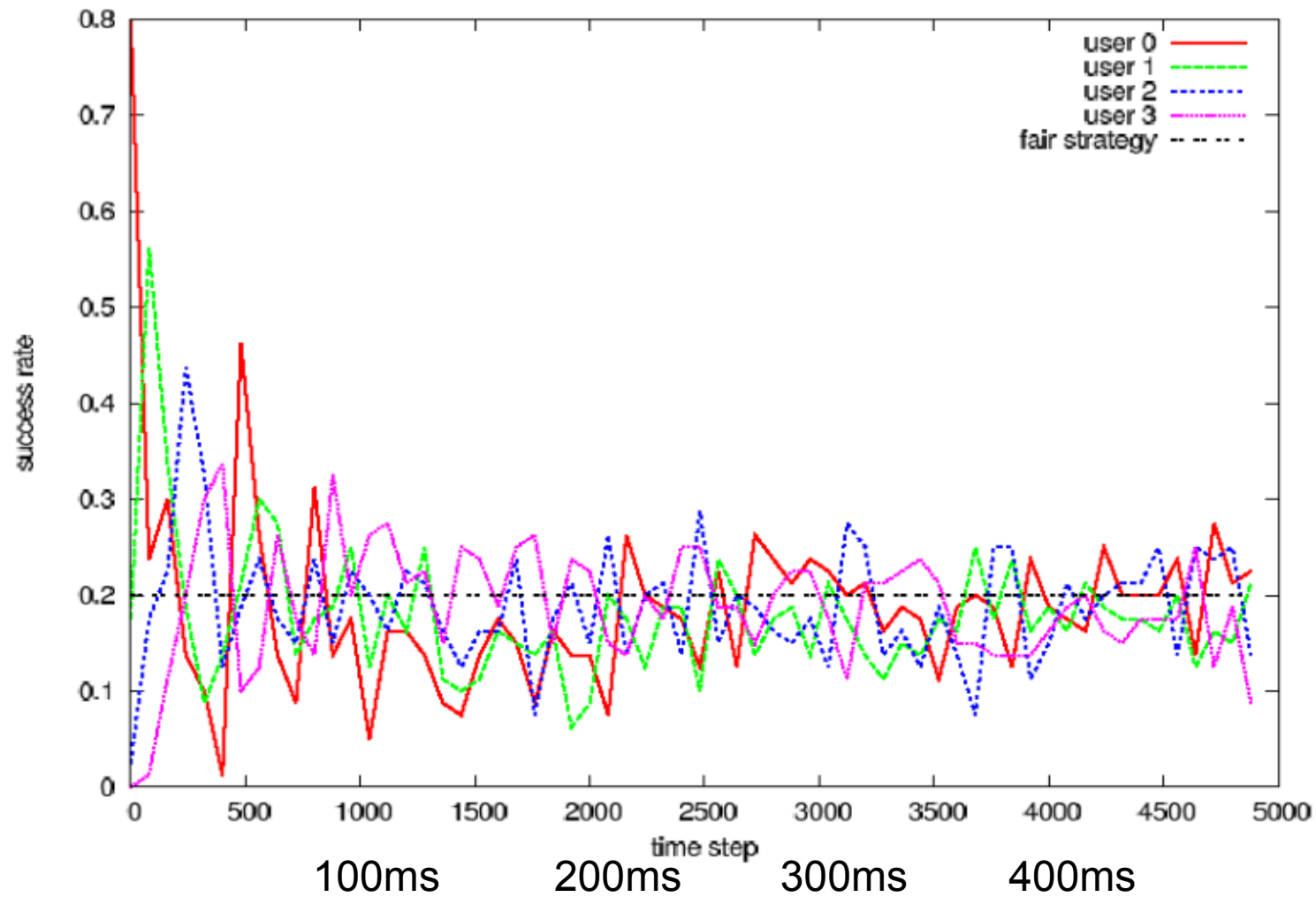
- It can be shown that a fair bandwidth sharing is one Nash Equilibrium of the enhanced game:

$$u_1(\mathbf{p}) = u_2(\mathbf{p}) = \dots = u_n(\mathbf{p}) = \frac{1 - \varepsilon}{n}$$

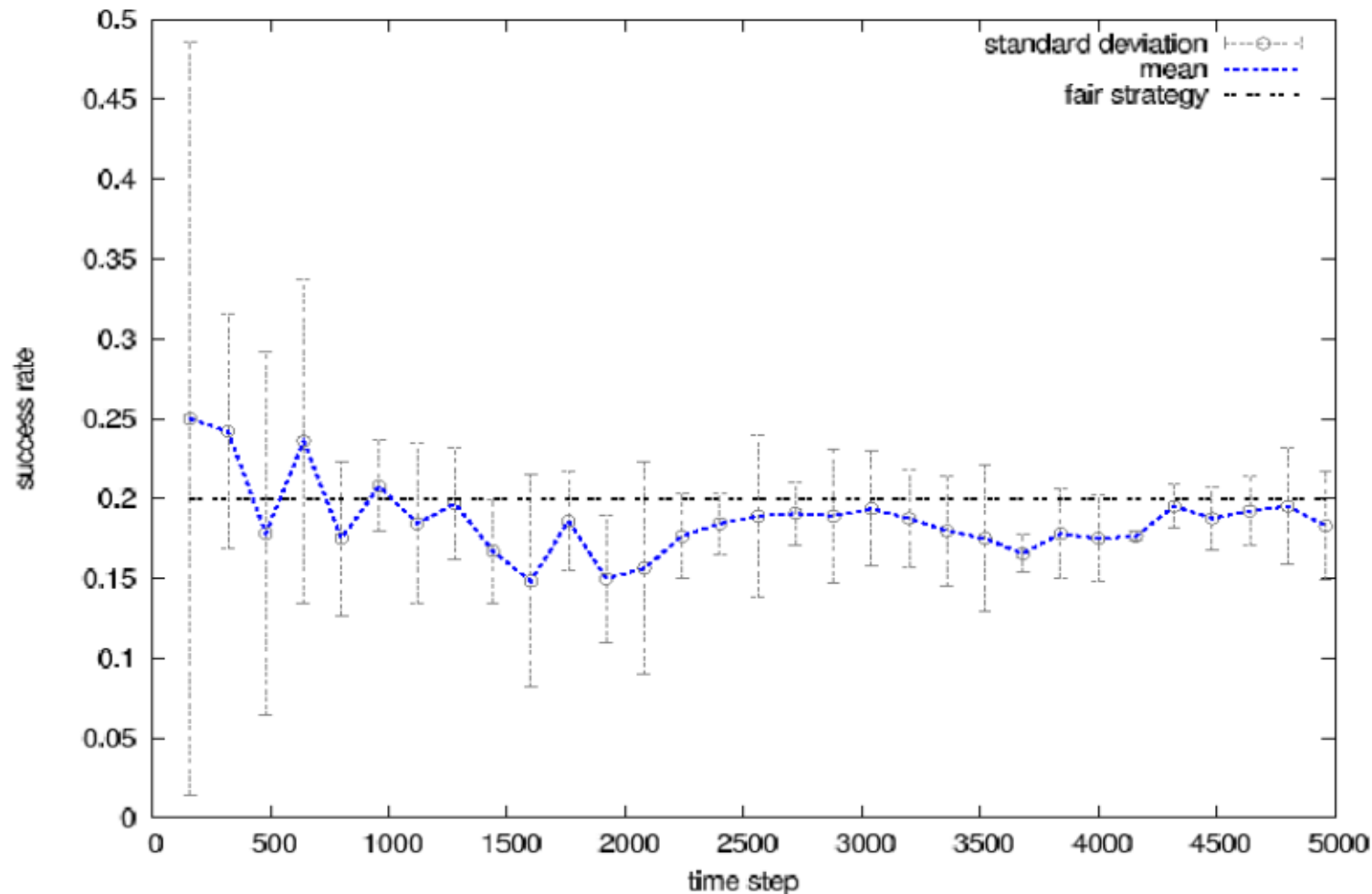
- This means, it is possible to design a self-organizing system for fair bandwidth sharing.
- However, the Nash Equilibrium is not a unique one.
- Providing a Multi-Agent Reinforcement Learning algorithm to lead the system to a fair strategy

Experimental Results – 4 Users

Success rates (= bandwidth used) of 4 Players and $\epsilon = 0.2$

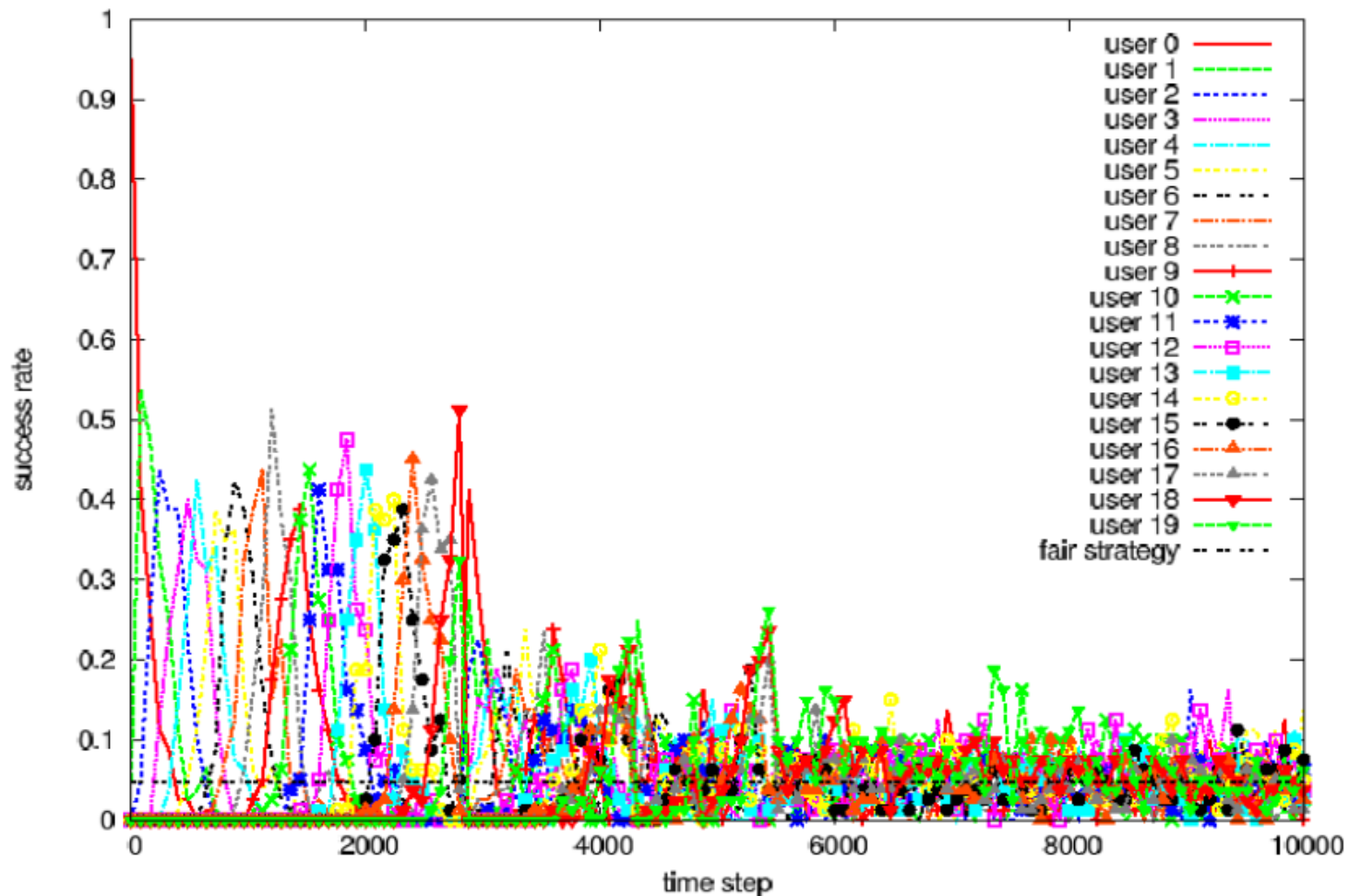


Experimental Results – 4 Users



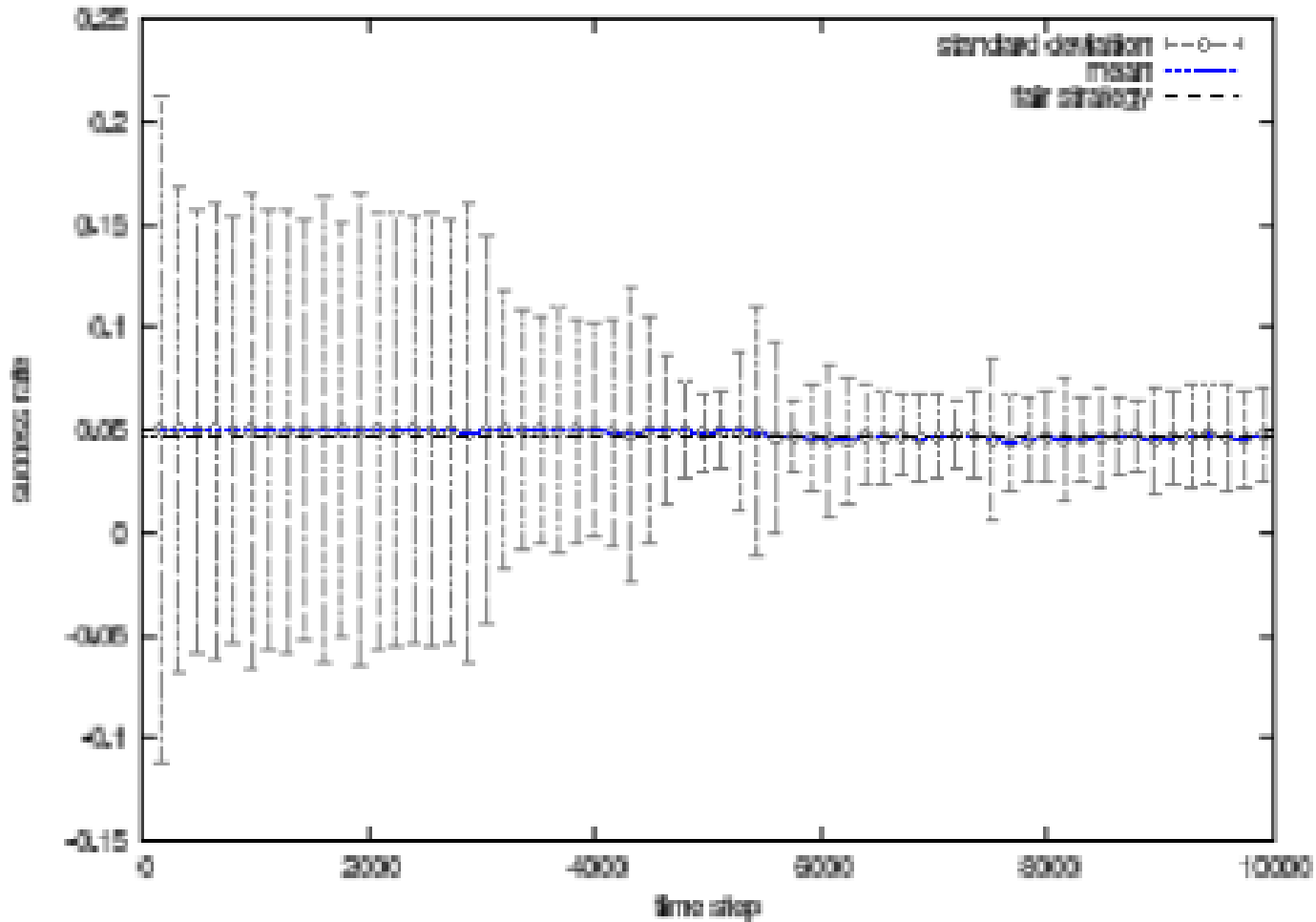
mean success rate with standard deviation

Experimental Results – 20 Users



20 Players and $\epsilon = 1/21$. Mean success rate and standard deviation.

Experimental Results – 20 Users



20 Players and $\epsilon = 1/21$. Mean success rate and standard deviation.

Roadmap

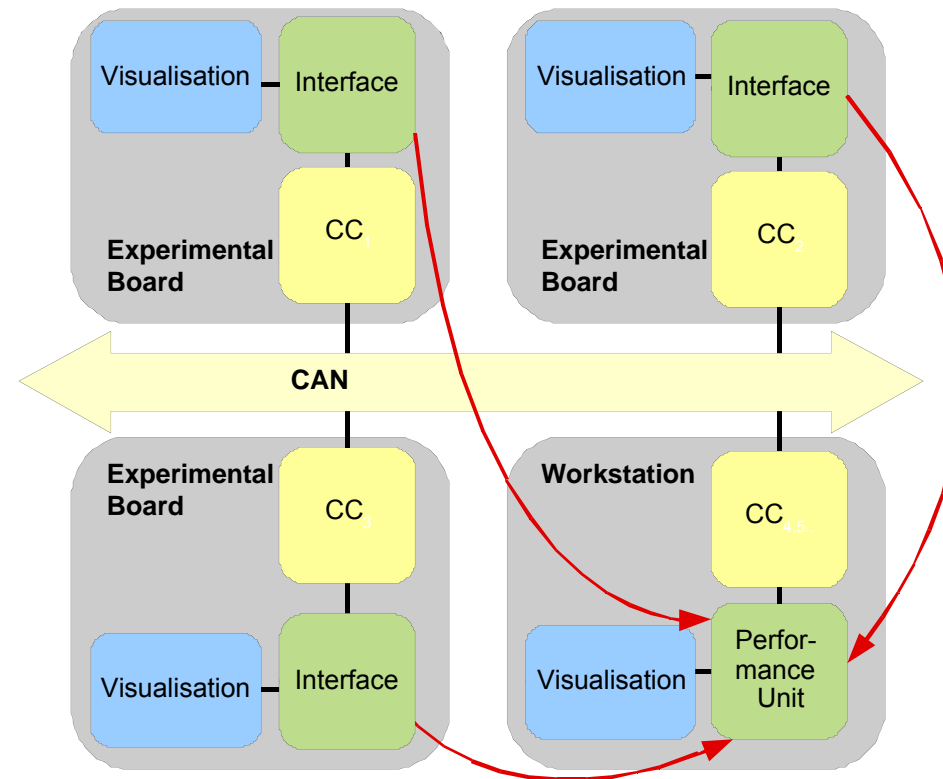
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Physical Simulation

- First step towards real implementation
- Analyze overhead of communication protocol
- Consideration of asynchronous communication with Controller Area Network (CAN)

Implementation on Real Hardware

- Implementation of organic bus arbitration algorithms on FPGA-boards
- Connecting boards and workstation via CAN
- Additional point-to-point connections for debugging purpose



Conclusions

- 3 Steps to achieve *Organic Self-organizing Bus-Based Communication Systems*:
 - Theoretical foundation (game theory)
 - Models and learning techniques
 - Real hardware implementation

Thanks for your Attention

- Project page:
 - www12.informatik.uni-erlangen.de/research/organicbus/
- Contact:
 - Tobias Ziermann
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Online Protocol for Adaptive Agents

- Collect local information by observing the environment, i.e., the shared medium, for some rounds.
 - *SuccessRate*: the rate of successfully arbitrating the medium
 - *ContentionRate*: the rate the medium was not free

- Update the sending probability using a learning algorithm:
Let Δ be the change of the probability
IF *ContentionRate* > 1 - ϵ **THEN**
 - // Coordination behavior behavior
 - $\Delta = -\textit{SuccessRate}$;**ELSE**
 - // Payoff maximization behavior
 - Apply weighted policy learning algorithm inspired by [2];**END IF**

References

- [1] Sudipta Rakshit; Guha, R.K., "**Fair bandwidth sharing in distributed systems: a game-theoretic approach,**" *Computers, IEEE Transactions on* , vol.54, no.11, pp. 1384-1393, Nov. 2005
- [2] Abdallah, S. and Lesser, V. 2006. "**Learning the task allocation game,**" *Proceedings of the Fifth international Joint Conference on Autonomous Agents and Multiagent Systems. AAMAS '06.* ACM, New York, NY, 850-857.